

Research in the Verification of Flight-Critical Systems

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Overview

- Background
- Verification and Certification
- Aircraft self-separation
- Runtime verification
- Formal-Methods for numerical software
- Experimental research

NASA Langley

- LaRC created in 1917 as the first National Advisory Committee for Aeronautics (NACA) research facility
 - Located in Hampton, Virginia
 - Primarily R&D focus
- LaRC became a NASA lab in 1958
 - The Mercury program began at LaRC
- Research areas of focus: Aeronautics, Atmospheric Sciences, and Exploration
 - Aerospace engineers dominate the research culture
- I belong to Safety-Critical Avionics Systems Branch

NASA R&D in Formal Methods

- NASA Langley Research Center (LARC) Safety Critical Avionics Systems Branch
 - Fault-tolerance
 - Air Traffic management
 - Theorem proving
- NASA Ames Research Center (ARC) Robust Software Engineering Group
 - Model Checking
 - Static analysis
- Jet Propulsion Laboratory (JPL) Laboratory for Reliable Software
 - Mission oriented
 - Mars rover software



VERIFICATION & CERTIFICATION

NASA

So, You Want to Build a Commercial Aircraft?

- Form a startup and start hacking just like silicon valley right?
 - Not so fast!
- Process starts off with a notification of intent to the FAA
 - A minuet begins between the company and the regulators
 - For a Part 25 aircraft they will tell you over 1500 safety criteria you must meet
 - Autos and medical devices are easy in comparison
 - DoD aircraft not subject to these regulations
- The Federal Aviation Administration (FAA) must certify the aircraft
 - Designated Engineering Representative (DER)
- The cyber-physical component is one of the largest risk factors
- Verification ≠ Certification
 - Safety is a systems level property



We are very good at building complex software systems that work 95% of the time---but, we do not know how to build complex software systems that are ultra-reliable and safe.

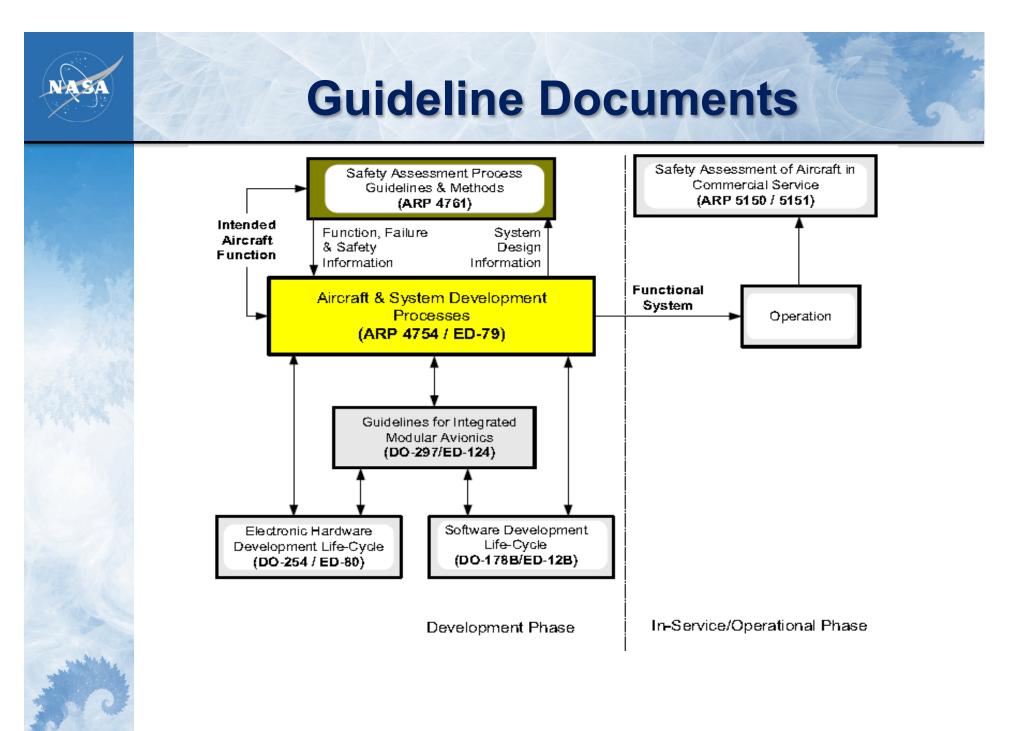
What, then has saved us in the past?

- -minimal amount of software that is safety critical
- -simple designs
- -Enormously-expensive verification and certification processes -backups that are not software, e.g.
 - ° hardware interlocks
 - ° human intervention

All sectors of aerospace are increasingly relying on software to perform safetycritical functions Size and Complexity

Software and Safety

- Critical avionics software is typically controlling some aspect of the aircraft
 - Control surfaces, fuel, ...
- System must continue to operate safely in the presence of hardware failures
 - Byzantine faults are a reality in this environment
- Systems must be engineered to be be safe
 - Human is a critical component
- Burden to handle the added complexity to ensure safety usually falls on the software and the humans in the cockpit



Eliminating Common Mode Errors

- Independence A concept to minimize the likelihood of common mode and cascade errors
- Diversity
 - HW, SW,
- Redundancy
 - Triple redundancy
 - Com/Mon
- Can mix techniques
 - Dissimilar com/mon

Formal Methods I

- Aerospace industry has used formal methods to analyze architectural properties of designs
 - TTE protocols
- Existing certification regime very process/test oriented
- DO-333 is an RTCA standard allowing formal methods to replace unit test for evidence
 - Standard explicitly mentions: syntax, semantics, soundness
 - Still resolving tool qualification questions

Formal Methods II

- Need much more work on languages and tools for specifying and analyzing architectures and designs of complex distributed hard real-time systems
 - Industry typically ignores the semantics until too late
- Avionics software much more restrictive than most commercial software
 - No recursion
 - No dynamic memory allocation allowed
 - Real-time scheduling issues always an issue
 - Lots of numerical code
- Most existing static analysis efforts not focused on this class of code
 - Very small market

Certification I

- Certification authorities certify an aircraft as a whole
 - You build everything in conformance to standards and processes
 - DERs sign off every step of the way
 - No provision for certification by composition of certified modules
- Why?
- Hint: Certification is about safety

Certification II

- Safety is not a compositional property
- Can assume/guarantee reasoning help?
 - Assumptions must include a fault model
 - How does system behave when assumptions fail
- Little work has been done in identifying the hurdles for modular certification
 - Rushby "Modular Certification"
- Challenge lies in the intersection of formal verification, fault tolerance, and safety-engineering



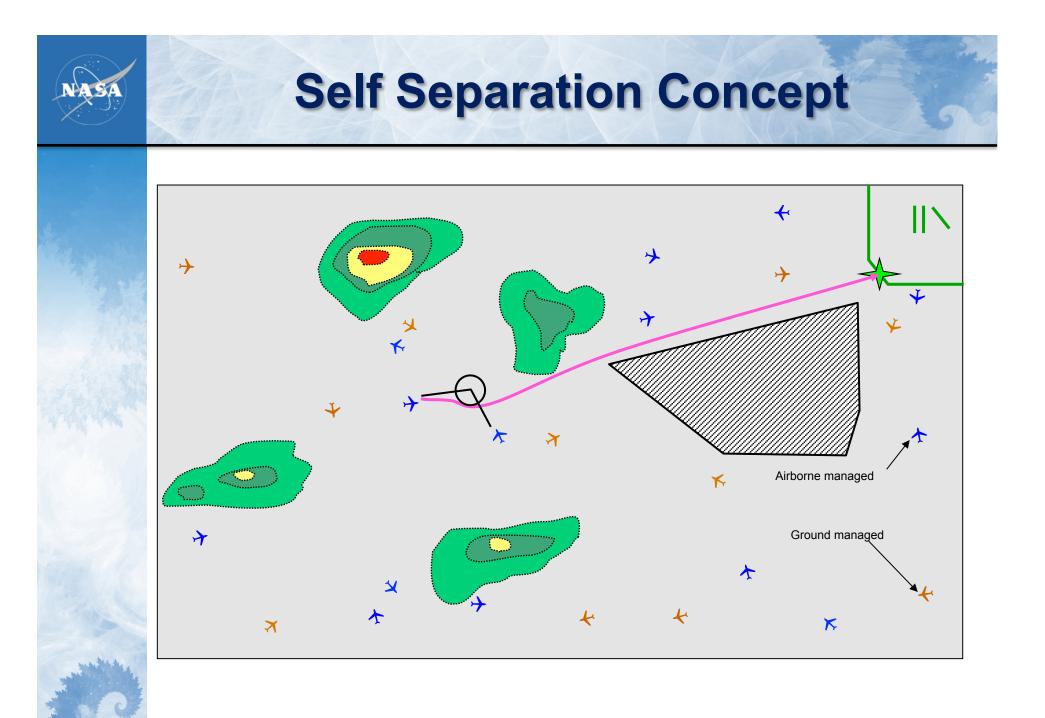
SELF SEPARATION CONCEPT

Aircraft Separation

- NASA is investigating a variety of air traffic management concepts to look at increasing capacity, efficiency, flexibility, etc.
- Adding more controllers will not achieve gains in these parameters
- Enabled by Automatic Dependent Surveillance Broadcast (ADS-B)
- Idea is to place more information in the hands of the pilots and trust them to make the right decision

Safe Self Separation

- More automation doesn't remove safety issues, but simply shifts the risk from people to automation
- Simulation helps find bugs
- Formal methods help show correctness
- Automated tools such as model checking and SMT solvers of little use due to lots of continuous math
- Interactive theorem proving is required

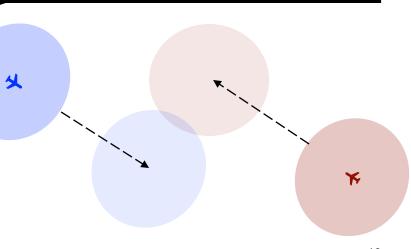


Separation and Automation

- Collision
 - Scrape paint



- Avoid through pilot, controller, and TCAS
- Loss of Separation
 - Separation standards are violated (5nmi, 1000ft)
 - Avoid through human and/or automation decisions
- Conflict
 - Predicted loss of separation



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Separation Algorithms

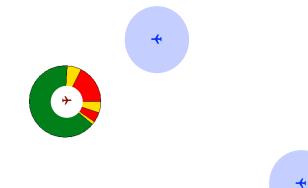
Conflict Detection

 Detect future loss of separation

Conflict Prevention – Provides ranges of conflict-free and conflict prone maneuvers

Conflict Resolution

 Suggest maneuvers to resolve a conflict





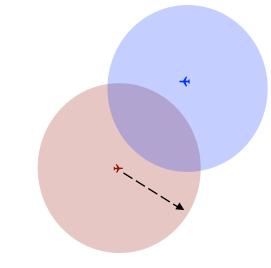
Recovery Algorithms

Conflict Recovery

 Suggest maneuvers to regain desired path

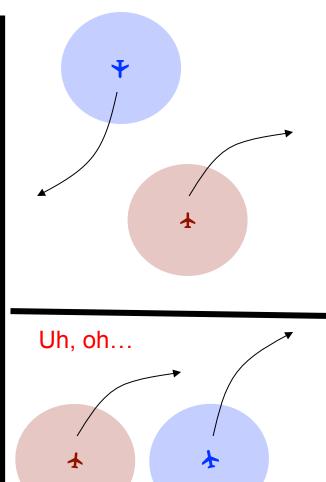
Loss of Separation Recovery

- For a variety of reasons separation may be lost
- Suggest a maneuver to regain separation



Conflict Resolution

- Each aircraft determines its own maneuvers through a combination of:
 - Go right/left, Speed up/slow down, Go up/down
- Properties
 - Independence: free of conflicts if one aircraft maneuvers
 - Coordination: free of conflicts if both aircraft maneuver
- Requirements
 - No specific comm between aircraft
 - No unfair rules: lower aircraft ID goes first, etc.



Formal Statement of Properties

independent: THEOREM

```
precondition_ind?(s(a), s(b), v(a), v(b)) AND
```

```
(nva = cr3d_vertical_speed(a,b) OR
```

```
nva = cr3d ground speed(a,b) OR
```

```
nva = cr3d heading(a,b)) AND
```

IMPLIES

NOT conflict?(s(a), s(b), nva-v(b))

coordinated: THEOREM

```
precondition_coord?(s(a), s(b), v(a), v(b)) AND
(nva = cr3d_vertical_speed(a,b) OR
nva = cr3d_ground_speed(a,b) OR
nva = cr3d_heading(a,b)) AND
(nvb = cr3d_vertical_speed(b,a) OR
nvb = cr3d_ground_speed(b,a) OR
nvb = cr3d_heading(b,a))
IMPLIES
NOT conflict?(s(a),s(b),nva-nvb)
```

ACCoRD Framework Airborne Coordinated Conflict Resolution and Detection ullet(ACCoRD) – a verification framework for *classes* of separation algorithms Safety Property Complex proof that criteria is safe -- provided by ACCoRD (Hopefully) straightforward Criteria proofs that each algorithm satisfies the criteria Algorithm A Algorithm B Algorithm C 24

Criteria is Very General

- The criteria was developed to aid the verification process
- Criteria allows combination of horizontal and vertical maneuvers
- But even more, if different algorithms satisfy the criteria, then they will be coordinated with each other
- Self separation does not rely on everyone running the same algorithm!

Criteria

| | in Conflict | in Loss of Separation |
|--------|---|--|
| horizo | ontal (s∙ v′) ≥ ε R (s [⊥] • v′) | s • v' > s • v AND s • v' ≥ s (D - S)/T _h |
| vertio | $\Delta > 0 \text{ AND } t > 0 \text{ AND}$ $\delta = 1 \text{ AND } s_z v_z \ge 0$ OR $ s_z + t v_z \ge H \text{ AND}$ $\delta (s_z + t v_z) v_z \le 0$ | $v_z' \neq 0 \text{ AND } s_z v_z' \ge 0 \text{ AND}$ $s_z v_z \ge 0 ==>$ if $v_z = 0$ then break_symm(s)(v_z') > 0 else $sign(v_z) v_z' \ge 0$ |



RUNTIME VERIFICATION

Autonomy Research for Civil Aviation: Toward a New Era of Flight

National Research Council Autonomy Research for Civil Aviation: Toward a New Era of Flight. Washington, DC: The National Academies Press, 2014 The committee did not individually prioritize these barriers. However, there is one critical, crosscutting challenge that must be overcome to unleash the full potential of advanced increasingly autonomous (IA) systems in civil aviation. This challenge may be described in terms of the question "How can we assure that advanced IA systems—especially those systems that rely on adaptive/nondeterministic software—will enhance rather than diminish the safety and reliability of the NAS?" There are four particularly challenging barriers that stand in the way of meeting this key challenge:

- Certification process
- •Decision making by adaptive/nondeterministic systems
- •Trust in adaptive/nondeterministic IA systems
- Verification and validation

Runtime Verification Motivation

- Given the current state-of-the-art not all code can be formally verified
 - Code base too large
 - Complex logic
- Learning algorithms are a particular challenge
- How do we ensure that assumptions that underline safety are actually correct
- Runtime Verification part of the solution

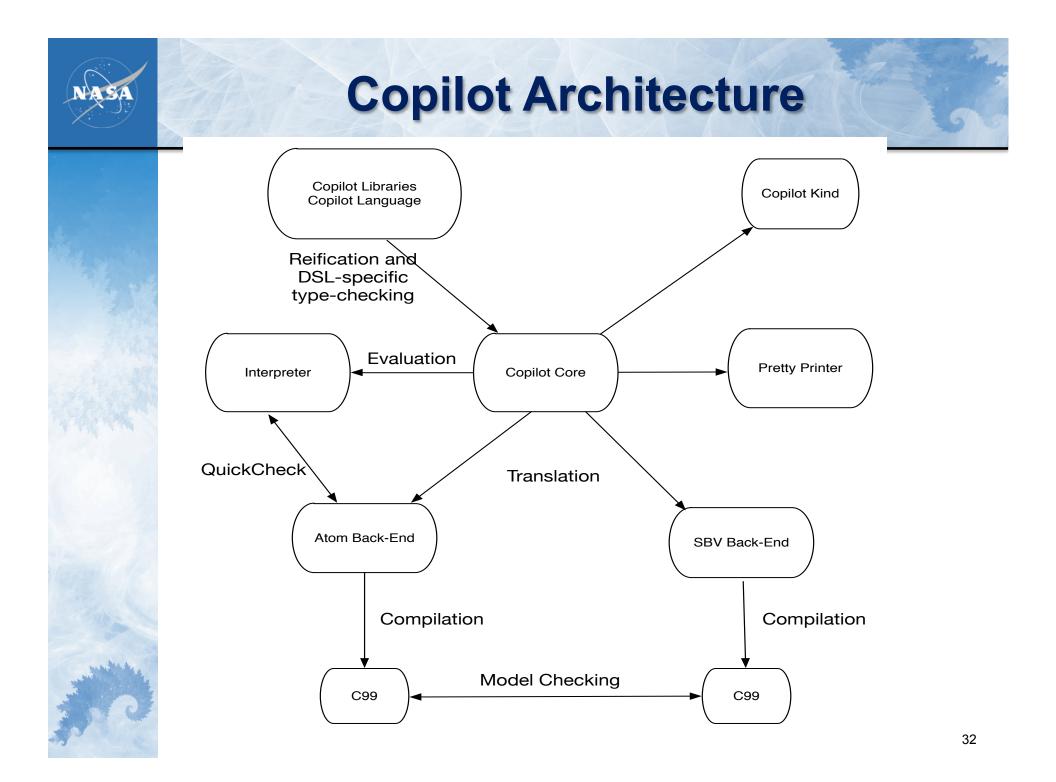
Runtime Verification

- System under observation (SUO)
- Correctness property φ
 - Past-time temporal logic
 - Regular languages
- Monitors observe SUO and detect violations of $\boldsymbol{\phi}$
- Accept all traces admitting φ

Runtime Verification in Copilot

- Copilot is a runtime verification framework aimed at hard real-time systems
 - Co-developed by Lee Pike at Galois and myself
 - Many Haskell hackers: Robin Moriset, Nis Wegmann, Sebastian Niller, Jonathan Laurent
- Written as a Haskell EDSL
- Composed of approximately 3000 lines of Haskell
- Copilot type system is embedded in Haskell's

 Hindley-Milner extended with type classes
- Translates into C99 through Atom or SBV



Copilot Language Operators

- Stream language constants and arithmetic operations are lifted to stream level
- (++) :: [a] -> Stream a -> Stream a
- xs ++ s prepends list xs to stream s
- drop :: Int -> Stream a -> Stream a
- drop ns skips the first n values in the stream
- Copilot specs must be causal stream values cannot depend on future values

Sample Copilot Specification

```
Haskell
```

```
fib :: [Word32]
fib = [0,1] ++ zipWith (+) (drop 1 fib)
```

```
Copilot
fib :: Stream Word32
fib = [0,1] ++ (fib + drop 1 fib)
```

Special constructs for input (sampling) and output (triggers)



Triggers

- Triggers provide a mechanism for Copilot streams to affect the outside world
- trigger:: String -> Steam Bool -> [TriggerArg] -> Spec
 - Name of external function
 - Guard determining when trigger is executed
 - List of arguments passed to the trigger

Trigger Example I

- If the temperature rises more than 2.3 degrees within 0.2 seconds, then the fuel injector should not be running
- Assuming that the global samplerate is 0.1 seconds

propTempRiseShutOff :: Spec propTempRiseShutOff = trigger "over_temp_rise" (overTempRise && running) []

Trigger Example II

where

max = 500 -- maximum engine temperature temps :: Stream Float temps = [max, max, max] ++ temp temp = extern "temp" Nothing overTempRise :: Stream Bool overTempRise = drop 2 temps > (2.3 + temps) running :: Stream Bool running = extern "running" Nothing

Watching the Watchers

- Lightweight verification techniques applied to ensure the generated code is safe and correct
 - Model checking
 - Quick check
- SMT-based model checking applied to verify monitor properties
 - U of Iowa's Kind2 new IC3 based model checker employed

NASA

Copilot Experiments

Tasks

- Use monitors with custom HW to bolt on fault tolerance
- Flight tests to validate concept

Edge



Future Research

- Flight Control software running on RTOS
 - Explore scheduling and architectures
 - Reachability analysis to derive monitors
- Simplex architectures that we can certify
- Monitor assumptions that underline safety cases
- Runtime verification applied to geofencing
- Using Copilot to verify autonomous software



VERIFICATION OF NUMERICAL SOFTWARE



From Models to Programs

- Our models and proofs on aircraft separation are carried out on real numbers in PVS
- Computers running on aircraft do not execute exact real arithmetic!
- A theorem that depends on exact mathematical analysis may fail to hold in the inexact realm of floating point numbers
- Recent effort to focus on these issues
 - A different perspective than our scientific computing and numerical analysis experts

Frama-C

- Frama-C is an OCAML-based platform for the static analysis of C programs developed at CEA
- Frama-C uses external decision procedures to prove ACSL annotations of C functions
- WP is a plug-in named for Dijkstra's Weakest Precondition calculus for deductive verification of programs
- Epsilon is an new project to add special functionality to WP to analyze numerical programs

Separation of Concerns

- Verify the functional correctness of air traffic management algorithms in PVS
 - Higer-Order Logic
 - Reasoning over R
- Implement the algorithm in C
- Annotate the code using ACSL
 - First-Order logic
- Use Frama-C deductive verification tool to prove numerical error bounds on operations preserve safety argument

Frama-C Specification + Code

/*@ logic real tauR(real ux, real uy, real vx, real vy, real b, real t) =
@ dmin(dmax(b*sqvR(vx,vy), - dotR(ux,uy, vx,vy)), t*sqvR(vx, vy)); @*/

/*@ requires -185200. <= s_x <= 185200. && -185200. <= s_y <= 185200. && @ -308. <= v_x <= 308. && -308. <= v_y <= 308. ; @ ensures \abs(\result - tauR(s_x, s_y, v_x, v_y, Br, Tr)) <= 0x1.p-8; @*/

Epsilon

- Represent floating point values as intervals
 - Applied interval analysis
- Represent numbers as finite sum of powers of 2
- Keep track of relative error and absolute error at each operation
- We want to use static analysis techniques to estimate the error
- Apply techniques from abstract interpretation to the logical specification



EXPERIMENTS

Edge 540T UAS Test Bed



Initial development funded by ARMD/IVHM now ARMD/SSAT

- Battery health prognosis
- Software health
- Autonomous decision system algorithms
- Current Capability:
 - Autopilot flight plan following
- Battery prognostics algorithms
- ADS-B IN/OUT
- Air traffic conflict detection and resolution
- Available Data:
 - Airframe parameters
 - Power plant parameters
 - INS/GPS
 - ADSB IN traffic table
- Advantages:
 - Low cost vehicle suitable for high risk avionics algorithm and hardware testing.
 - Low cost air traffic scenario flight testing.
 - Minimal ground support hardware requirements.
 - Environment for testing decision algorithms for space applications.



FAA registered LaRC 8' UAS: "N802RE" Soon to be registered: "N803RE"

~ \$20k per copy

Test Concept

